

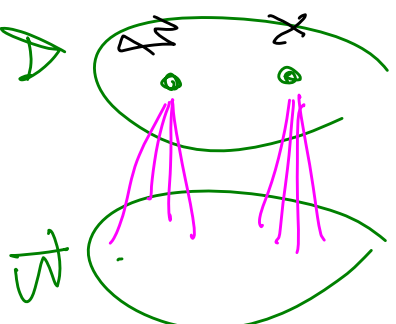
E.g. show (must be true);

$r=3$ | a.s. \nexists max cut (A, B) s.t.

$x, y \in A$ \bar{w} $d_B(x, y) = 0$

$\phi > Cn^{1/2} \sqrt{\log n}$

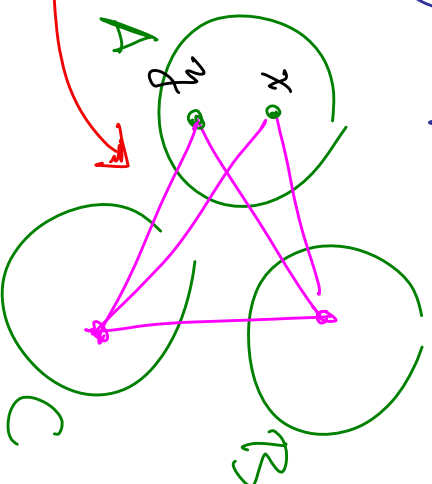
Should be $\approx np^2/2$



$r=4$ | a.s. \nexists max "cut" (A, B, C) s.t.

$x, y \in A$ \bar{w} $K(x, y, B, C) = 0$

number of these



(should be $\approx n^2 p^5 / q$)

$$\text{I a.s. } \Pi = (A, B) \text{ M.V.C. } \Big| \Rightarrow d_B(x, y) \neq 0 \quad (*)$$

$x, y \in A$

PF (-1) a.s. $d(v) \approx np \forall v$ (~~SWMA~~)

Fix x, y

$$(3) (A, B) \text{ M.V.C. } \Big| \Rightarrow d_B(x) \geq d(x)/2$$

$x \in A$

$(\approx np^2/2)$

\rightarrow ETS ~~(*)~~ for

$$\Pi \in \mathcal{G} = \{ (A, B) : x, y \in A, \underline{d_B(x)} \approx np/2 \}$$

Pr (1) Choose G :

(a) $N(x)$

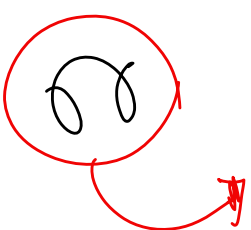
(b) rest of E except $\nabla(y, N(x))$ } G'

(c) $\nabla(y, N(x))$

Pr $\nabla^* \equiv (S, T)$:

\max^* among cuts of G' in G

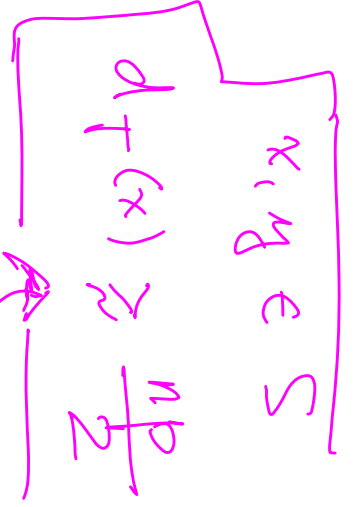
$[x, y \in S$
 $d_T(x) \geq np/2$



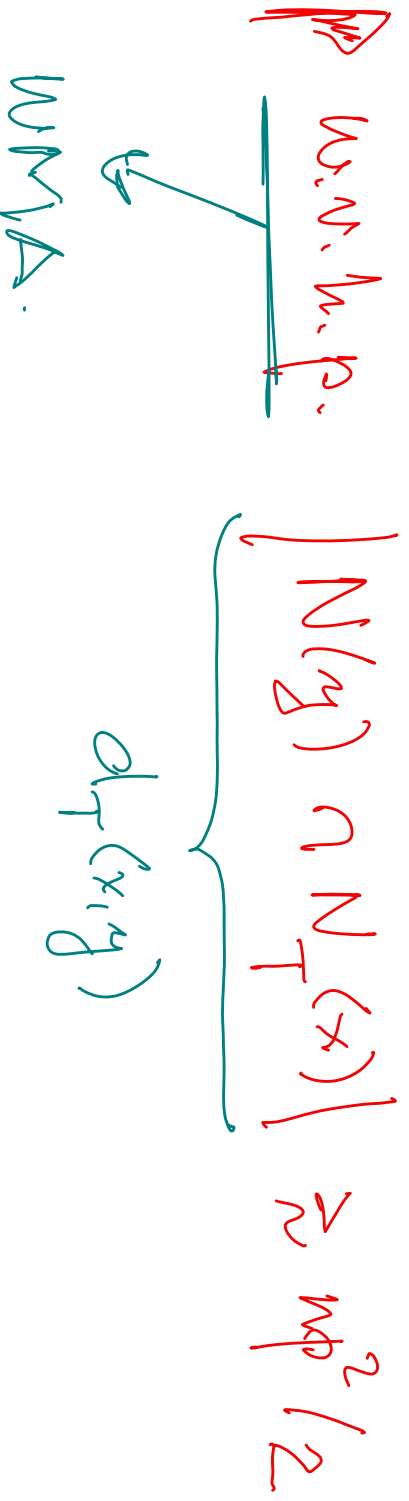
* first max ...

- (a) $N(x)$
- (b) rest except $\nabla(y, N(x))$
- (c) $\nabla(y, N(x))$

$\Pi^* = (S, T) : \max$ for G' among cuts in \mathcal{C}



$$p > c \left\| \frac{cgn}{n} \right\| \implies \text{in } \underline{G} :$$



$$\text{I} \quad \text{In } \mathcal{G}' : \Pi^* = (S, T) : \max \text{ in } \boxed{\mathcal{E}}$$

$$\text{in } \mathcal{G} : \underline{\underline{d_T(x, y) \approx np^2/2}}$$

$$\text{In } \mathcal{G} : \Pi = (A, B) \in \mathcal{E} :$$

$$|\Pi_{\mathcal{G}}| = |\Pi_{\mathcal{G}'}| + d_B(x, y)$$

$$|\Pi_{\mathcal{G}}^*| = |\Pi_{\mathcal{G}'}^*| + \boxed{d_T(x, y)} \approx np^2/2$$

$$\text{M.C.} \implies d_B(x, y) \geq d_T(x, y) (\approx np^2/2)$$

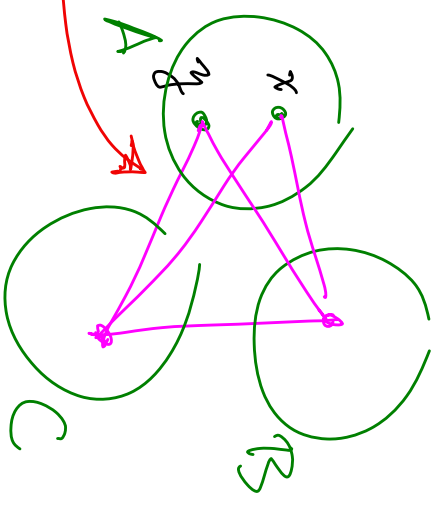
(more than we asked)

Q.E.D.

$(r=d)$ a.s. $\nexists (A, B, C) \text{ Mc } \nexists x, y \in A$

$$\overline{w} \quad \underline{K(x, y, B, C)} = 0$$

number of these

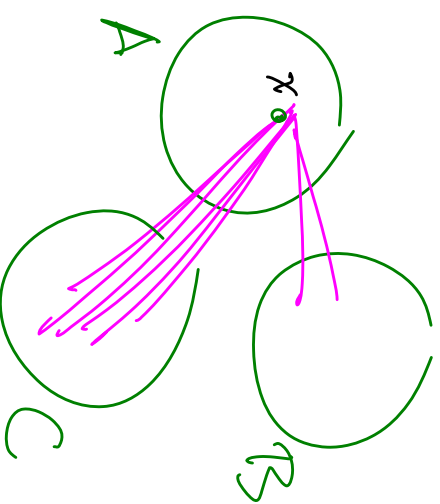


whn ... how about:

a.s. $\nexists (A, B, C) \text{ Mc } \nexists x \in A \quad \overline{w}$

~~*~~ $d_B(x) = \begin{cases} 0 \\ \text{small} \end{cases}$

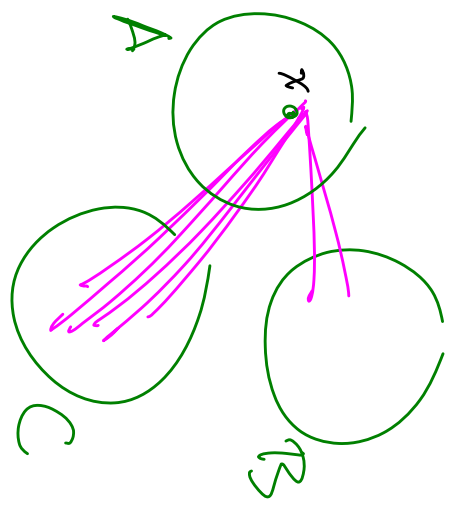
??



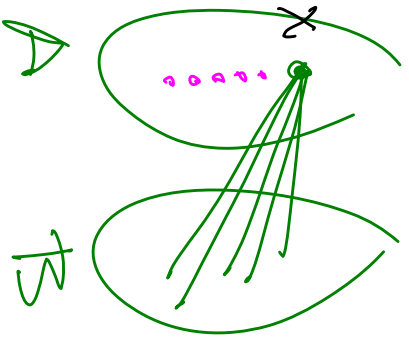
Γ a.s. $\nexists (A, B, C)$ MC $\forall x \in A$ \bar{w}

$\textcircled{*} d_B(x) = \begin{cases} 0 \\ \text{small} \end{cases}$

? ?



BTW don't know ($r=3$)

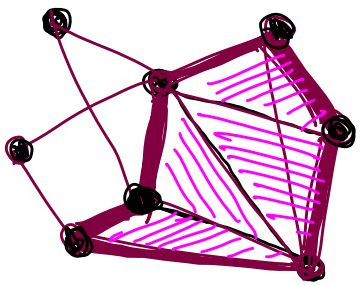


a.s. \nexists MC (A, B) $\forall x \in A$

$\bar{w} d_A(x) = 0$

3rd "group" (briefly): a topological g .

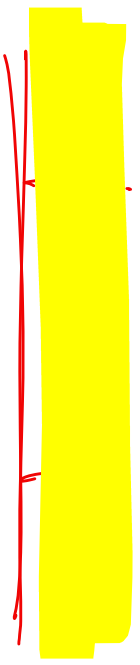
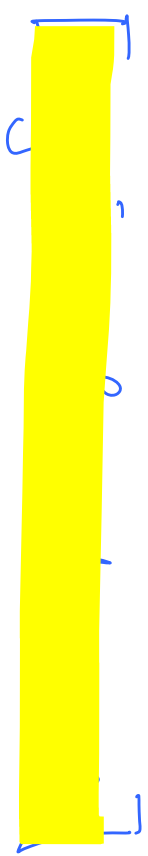
$\mathcal{F} = \{ \text{every cycle} = \text{mod } 2 \text{ sum of } \Delta\text{'s} \}$



a.k.a.

$$H_1(\Delta(G), \mathbb{Z}_2) = 0$$

\mathcal{F} cliques of G

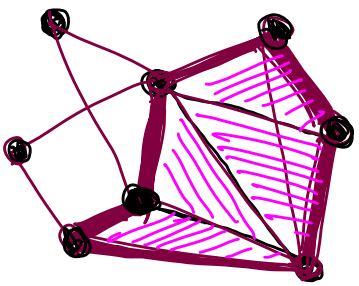


D-Hamm-K (Conj of M. Kahle):

$$p > (1 + \epsilon) \sqrt{\frac{3}{2} \frac{kmn}{n}}$$

$$\Leftrightarrow G \models \mathcal{F} \text{ a.s.}$$

$\mathcal{F} = \{ \text{every cycle} = \text{mod } 2 \text{ sum of } \Delta_i \}$



$$\underline{\text{Thm}} \quad p > (1+\epsilon) \sqrt{\frac{3}{2} \frac{2m}{n}}$$

$\Rightarrow G \models \mathcal{F}$ a.s. \downarrow

Real statement:

$\mathcal{Q}_a = \{ \text{every edge in a } \Delta_i \}$

$$\underline{\underline{\text{Thm}}} \quad Pr(G_{n,p} \models \mathcal{Q} \wedge \mathcal{F}) = o(1)$$

Conj (M. Kahle) $p > (1+\varepsilon) p_c(k, n)^*$

$$\implies \text{w.h.p. } H_k(\Delta(G), \mathbb{R}) = 0$$

anything

* $p_c(k, n) \leftrightarrow$ every K_{k+1} in a K_{k+2}

Kahle: true if $R = \mathbb{Q}$

(DHK: $k=1, R = \mathbb{Z}_2$)

all other cases open

Back to

general increasing $f \subseteq 2^X$

and "GAP"

(what drives p_e ?)

(pre-) Conj (K-Kalai OT)

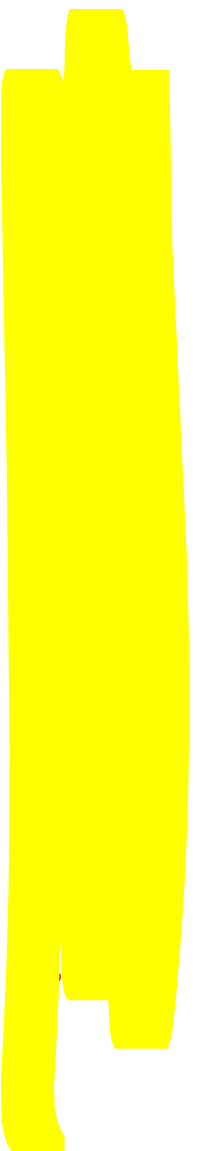
$\forall f$

\exists trivial d.b. $g(f)$ on $P_c(f)$ s.t.

$$P_c(f) < K g(f) \cdot \log |X|$$

\rightarrow const (e.g. $K=1$?)

$f \subseteq Z^X$
incr.



e.g. $G_{n,p}$, $\exists = \exists_H$ ($H = \text{something}$)

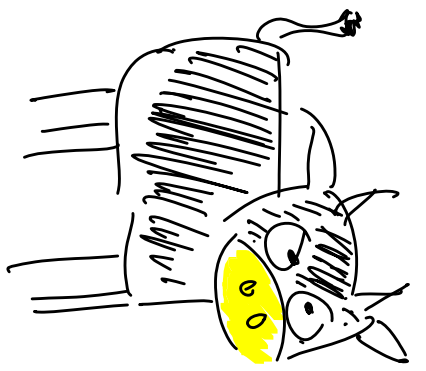
con' \rightarrow

if $\mathbb{E}_p[\#F's] > 1$ $\forall F \subseteq H$
 then $p_c(\exists_H) < k p \log n$

— and if

$\exists \neq \exists_H ?$

— con': also

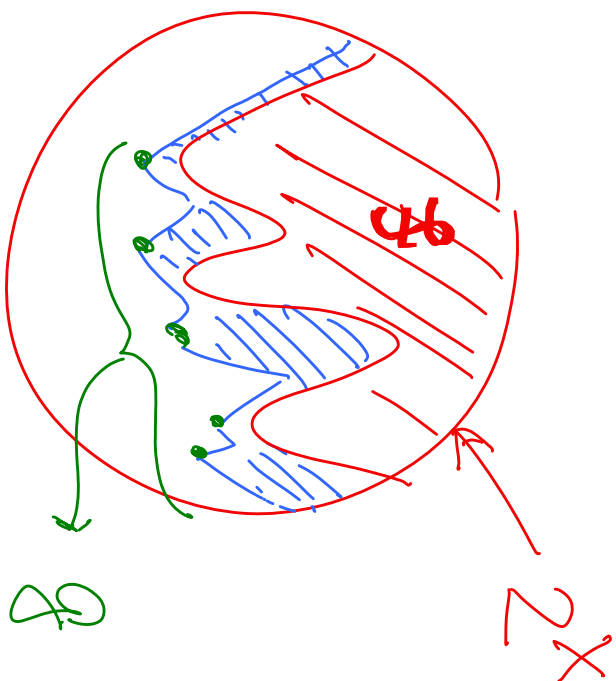


trivial proof that

$$P_c(\mathcal{F}) > P$$

$$\exists \mathcal{G} \subseteq 2^X \quad \mathcal{F}$$

$$\textcircled{1} \quad \left[\begin{array}{l} A, B \in \mathcal{F} \\ \exists A \in \mathcal{G}, A \subseteq B \end{array} \right]$$



$$\mathcal{F} \subseteq \langle \mathcal{G} \rangle = \{ B \subseteq X : \exists A \in \mathcal{G}, A \subseteq B \}$$

$$\exists A \in \mathcal{G}, A \subseteq B$$

$$\textcircled{2} \quad \sum_{A \in \mathcal{G}} P^{|A|} < \frac{1}{2}$$

unimp.

Triv. pf \otimes of $\Pr(\mathcal{F}) > p$:

$$\exists \mathcal{G} \subseteq \mathcal{Z}^X \text{ s.t. } \langle \mathcal{G} \rangle \supseteq \mathcal{F}$$

$$\textcircled{2} \sum_{A \in \mathcal{G}} p^{|A|} < \frac{1}{2}$$

\otimes

$$Y \sim \mu_p$$

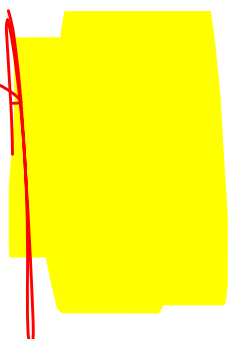
\implies

$$\frac{1}{2} > \sum_{A \in \mathcal{G}} p^{|A|} = \mathbb{E}_p \left[\sum_{A \in \mathcal{G}} \mathbb{1}_{A \subseteq Y} \right]$$

$$\Pr(Y \supseteq A)$$

$$\geq \mu_p(\mathcal{F})$$





$$:= \sup \{ p : \exists Q \} \leq p_c(f)$$

"expectation threshold"

triv.

CONJS (k-Kalai '07) $\exists k$

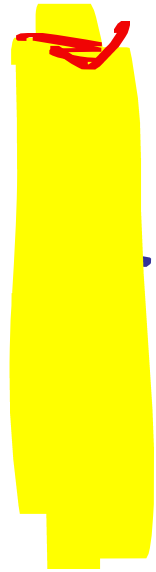
$\forall f$

$$p_c(f) < k p_E(f) \log |X|$$

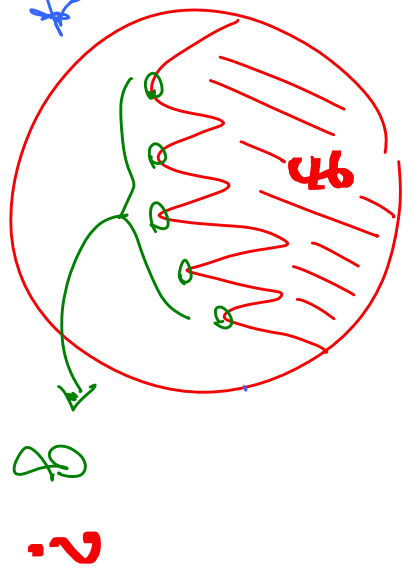
\equiv { tree conjs (EX) \longleftrightarrow show
 Shamir (easier EX) $p_E \leq \dots$

$$P_{\mathbb{H}}(\mathcal{F}) = \sup \{ \rho : \exists Q \ni \left\{ \langle Q \rangle \geq \mathcal{F} \right. \\ \left. \sum_{A \in Q} \rho^{|A|} < \frac{1}{2} \right\} \}$$

Can't $\rho_c < k P_{\mathbb{H}}(\log |X|)$



E.g. $Q = \min(\mathcal{F})$?



$$Q = \min(\mathcal{F}) = \{ H's \} \text{ BAD}$$

$$(Q = \{ \text{diamond 's'} \} \text{ GOOD})$$

E.g. Ramsey (R's)

$$\text{(recall } r_c = \Theta(n^{-1/2})$$

SHOW: $p < c n^{-1/2} \Rightarrow \exists G$ [so no gap]

we may regard pfs* as:

coloring algorithm \rightarrow failure finds G

* Wezale-Rucinski-Vaigk,

Rödl-Ruciński, Nenadov-Siegel

[coloring algorithm finds δ]

repeat until can't: remove edge on ≤ 1 

→ "core" S must contain one of z

① large " Δ -can" piece 

(unlikely: $\#$ "children" $\approx 3np^2$)

② H of size $O(k)$ \bar{w} $e_H > 2\sqrt{H}$

unlikely

(nextiv) $e_H \leq 2\sqrt{H} \Rightarrow H \rightarrow (k_3)_2$



But in general "what's Q ?"

— or "how do we find Q " —
is ??

(more on this soon)

Takegrand

another triv. pf. ^{*} of $P_c(\mathcal{F}) > p$:

$$\exists \alpha \in \mathbb{R}^X \rightarrow \mathbb{R}^+ \ni$$

$$\textcircled{1} \text{ BCF} \Rightarrow \sum_{S \in \mathcal{B}} \alpha_S \geq 1$$

$$\left(\longleftrightarrow \langle \mathcal{G}, \mathcal{F} \rangle \right)$$

$$\textcircled{2} \sum_S \alpha_S p|S| < \frac{1}{2}$$

* EX

$$\textcircled{1} \text{ B E F} \Rightarrow \sum_{S \in \mathcal{B}} \alpha_S \geq 1$$

$$\textcircled{2} \sum_S \alpha_S p(|S|) < \frac{1}{2}$$

names : f is

p -small if $\exists \underline{\underline{\delta}} \dots$

weakly- p -small (wpps) if $\exists \underline{\underline{\alpha}}$

(TRIV : p -small \Rightarrow wpps)

① wps_0 : LP relaxation of p -Small:

$$\boxed{wps_0} \quad \min \sum x_s p^{|s|}$$

subject to $x_s \geq 0 \quad \forall s$

$$\sum_{S \in A} x_S \geq 1 \quad \forall A \in \mathcal{F}$$

$$\boxed{p\text{-Small}} \quad \text{same w } x_S \in \{0,1\}$$

$$p_E^* := \sup \{ p : \exists \alpha \}$$

$$p_E \leq p_E^* \leq p_c$$

$$\underline{\text{Conj}}^* \quad \forall \varepsilon \quad p_c(\varepsilon) < k p_E^*(\varepsilon) \log |\lambda|$$

of course $\text{Conj } k k \Rightarrow \text{Conj}^*$,

but maybe equiv:

Conj 2 $\exists k \exists$

$$p_E(\varepsilon) > p_E^*(\varepsilon)/k \quad \forall \varepsilon$$

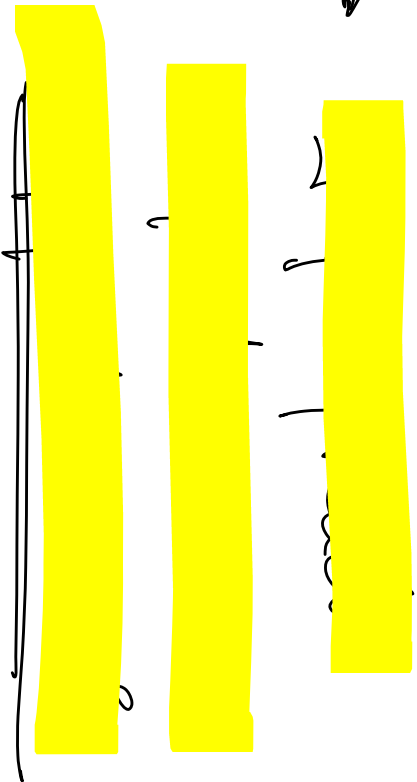
Big (?) advantage of Con^t : quality

Def $n: \mathbb{Z}^X \rightarrow \mathbb{R}^+$ is p-spread if

$$n(KIY) \leq p |I| \quad \forall I \subseteq X$$

$$\sum_{A \supseteq I} n_A = \mathbb{P}_r(A \supseteq I) \quad \text{if } A \sim n$$

$$\mathbb{P}_r^* \leq p \quad \longleftrightarrow \text{ESS.}$$



$\rho_E^* < p \iff \overset{\text{ess.}}{\exists} p\text{-sp. prob. meas. on } \mathcal{F}$
 TRIVIAL DIRECTION:

N p -sp. prob. m. on \mathcal{F}
 α as in wps

$$1 = \sum_{A \in \mathcal{F}} \nu_A \leq \sum_{A \in \mathcal{F}} \nu_A \sum_{S \subseteq A} \alpha_S \leq \sum_S \alpha_S p^{|S|}$$

$\implies \rho_E^* < p \quad \square$

~~Conj*~~ (again) If \mathcal{F} supports a p -spread
 prob. meas. then

Rank Usually in threshold problems we try to show $P_c < \text{natural guess}$.

For $k < K$ Conj. this is problematic:

want: $P_c \leq g \implies P_c < k g \log |x|$

[nonexistence of g for $g' > g$]

but don't know how to use hypothesis

\rightarrow start w P_c & try to show

$P_c \leq P_c / \log |x| \implies \exists g$ [more below]

OTM Conj. start w p-spread of (as above)

triv.: WMA $\text{supp}(\nu) \subseteq \mathcal{X} \stackrel{\circ}{=} \min \mathcal{F}$

E.g. if \mathcal{H} symmetric (Aut(\mathcal{H}) trans. on \mathcal{X})

$(\Rightarrow \mathcal{H} \subseteq \binom{X}{m})$ same m)

then (EX) WMA $\nu = \underline{\text{unif. meas.}}$ on \mathcal{X}

e.g.
EX \rightarrow ① $\nu(\langle x \rangle) = \frac{m}{|X|} \quad \forall x \in X$

② $\nu(\langle A \rangle) = \nu_A = \frac{1}{|\mathcal{X}|} \quad \forall A \in \mathcal{X}$

$\Rightarrow |\mathcal{H}| \geq \phi^{-m}$

E.g. Stochastic $\mathbb{P}(X = \binom{[n]}{3}) \quad \binom{3}{n}$

$\mathcal{X} = \{p, m, s\}, m = n/3, |\mathcal{X}| \approx \left(\frac{n^2}{2e^2}\right)^m$
EX for now

γ is unif. meas. on \mathcal{R} , $A \sim \gamma$

$$\left(\rightarrow \mathbb{P}(A \subseteq S) = \gamma(\langle S \rangle)\right)$$

$$\begin{aligned} \mathbb{P}(x \in A) &= \frac{m}{|X|} \approx \frac{2}{n^2} \\ \mathbb{P}(S \subseteq A) &\leq \left(\frac{2e^2}{n^2}\right)^{|S|} \quad A \subseteq S \end{aligned}$$

EX

[Any unif. meas. on \mathcal{X} ($= \{p, m, s\}$)

$$\Pr(x \in A) = \frac{m}{|X|} \approx \frac{2}{n^2}$$

$$\Pr(S \subseteq A) \leq \left(\frac{2e^2}{n^2}\right)^{|S|}$$

$$\approx \prod_{x \in S} \Pr(x \in A)$$

$x \in S$

\Rightarrow Coni* says this should be all we

need for $p_c(\mathcal{F}) < K \frac{\log n}{n^2}$ (as in SKV)

[still Talagrand, aiming for something stronger]

recall (2): $\mu, \lambda \in \mathcal{M}(Z^X) := \{ \text{prob. meas's} \}$
on Z^X ?

μ stochastically dominates λ ($\mu \succ \lambda$)

if $\mu(\mathcal{E}) \geq \lambda(\mathcal{E}) \quad \forall \text{ inc. } \mathcal{E}$

$\mu \neq \lambda: \mu(\mathcal{F}) \geq \lambda(\mathcal{F}) \quad \forall \text{ int. } \mathcal{F}$

$$\nu \text{ p-spread} \equiv [\mu_p(\langle S \rangle) \geq \nu(\langle S \rangle) \quad \forall S]$$

maybe $\mu_p \neq \nu$?

→ NO (stupid):

e.g. $\nu(\Sigma \neq \emptyset) = 0 < \mu_p(\Sigma \neq \emptyset)$

e.g. $\nu(\mathcal{F}) = 1 < \mu_p(\mathcal{F})$
supp'd on \mathcal{F}

BUT:

Def $q \in \mathcal{M}(Z^X)$, $\alpha \in [0, 1] \rightarrow$

$\lambda_\alpha(N) \in \mathcal{M}(Z^X)$: law of r.v. Z :

$$Y \sim q, Z = Y_\alpha$$

recall...

Conj A $\exists \alpha > 0 \exists$

$$N \text{ p-spread} \Rightarrow \lambda_\alpha(N) \leq \mu_p$$

⊙ would have important consequences

⊙ $\text{Conj A} \stackrel{\text{EX}}{\subseteq} \text{Conj B}$

Conj B $\exists L \exists$

ν g -spread $\dot{=} \alpha = 1 - (1-p)^{\frac{1}{L}}$

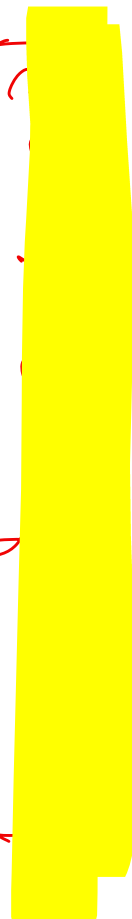
$$\implies \lambda_\alpha(\nu) \leq \mu_p \rightarrow$$

① ν g -spread on \mathcal{F} ,

$p = \frac{1}{g} \ln |X|$, α as above

Conj B \rightarrow

$$\mu_p(\mathcal{F}) \geq \lambda_\alpha(\nu)(\mathcal{F}) \geq \alpha |X| \geq 1/2$$

\rightarrow 



better:

$$\|Z\| := \max \{ |A| : A \in \min Z \}$$

above arg. gives (assuming Conj. B)

$$p_c(Z) \leq L_f \ln \|Z\| \rightarrow$$

$$\underline{\text{Conj}} \quad A \subseteq p_c(Z) < K p_c(\log \|Z\|)$$

[Rank: true if $\|Z\| = O(1)$]

(pf similar to Lyons)

back to Conj 2 ($\phi_E > \rho_E^* / k$)

[just an example for flavor]

G graph on X , $|G| = m$ ($p = \text{something}$)

Note $|E(G[X_p])| = mp^2$

$\exists W \subseteq X : |G[W]| \geq 200 mp^2$
(say)

G graph on X , $|G| = m$, $|E(G[X_p])| = mp^2$

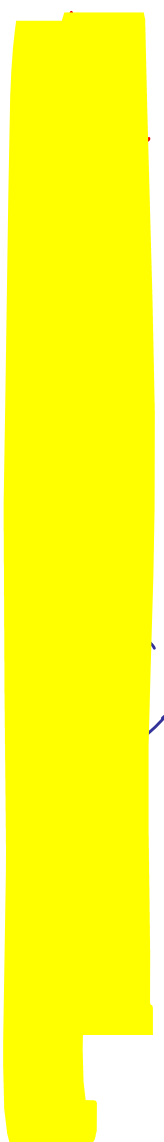
$$\mathcal{F} = \{W \subseteq X : |G[W]| \geq 200mp^2\}$$

ORS: $\chi_{\mathcal{F}}^*(\mathcal{F}) \geq 10p$:

$$\alpha_S = \frac{1}{200mp^2} \quad \forall S \in G \left(S \subseteq \binom{X}{2} \right)$$

$$\rightarrow \textcircled{1} W \in \mathcal{F} \Rightarrow \sum_{S \subseteq W} \alpha_S \geq 1$$

$$\textcircled{2} \sum \alpha_S (10p)^{|S|} = 1/2$$



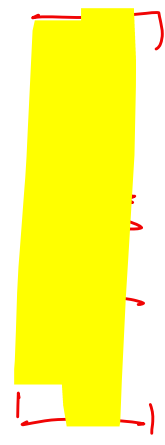
$i.e. \rightarrow$

$$\mathcal{F} = \{W \subseteq X : |G[W]| \geq 200 m p^2 \}$$

SHOW: $\exists q \subseteq \mathcal{F} \Rightarrow$

① $\langle q \rangle \in \mathcal{F}$

② $\sum_{A \in q} p^{|A|} < 1/2$



E.g. $|W| = 10np \Rightarrow W \in \mathcal{F}$ typical

but $q \ni \{ \text{such } W \}'s$ is small:

contrib $\approx \binom{n}{10np} p^{10np} < \left(\frac{e^{np}}{10np} \right)^{10np} = \text{tiny}$
 (to ②)

$\mathcal{F} = \{W \subseteq X : |G[W]| \geq 200 m p^2 \epsilon\}$
more interesting: atypically dense W 's

E.g. G d -reg, $W \in \mathcal{F}$, $|W| < 10np$

EX $\Rightarrow \exists x \in W : |N_x \cap W| > 10dnp$
 \leftarrow *nbhd*

\rightarrow add to \mathcal{G} :

all sets $\{x \in X \cup Y \mid \bar{x} \in Y \in \binom{N_x}{10dnp}\}$

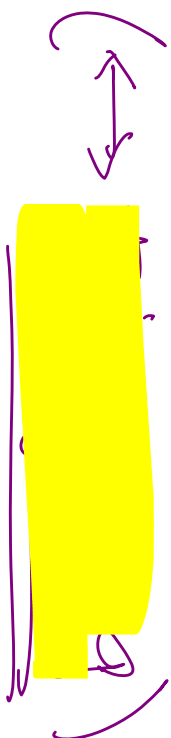
\rightarrow contrib: $n \binom{d}{10dnp} p^{10dnp+1} < n(e/10)^{10dnp}$

\rightarrow fine if $dnp \geq \log n$ (if not? ...)

[back to KK soon]

back to: how sharp is the threshold?

defs:



$$P_{\alpha}(z) = P(\bar{w} \leq \mu_{\bar{w}}(z) = \alpha)$$

e.g. $P_C = P_{1/2}$

"threshold interval": $\Theta(z) = [P_{0.1}, P_{0.9}]$

$$\delta = \delta z = | \Theta(z) | = P_{0.9} - P_{0.1}$$

$$\Gamma \delta = \delta f = p_{0,1} - p_{1,1} \Gamma$$

$$\text{th. is } \begin{cases} \underline{\text{COARSE}} & \text{if } \delta / p_c^* = \Omega(1) \\ \underline{\text{SHARP}} & \text{if } \delta / p_c^* = o(1) \end{cases}$$

$$* \text{ really } \delta(f_n) / p_c(f_n)$$

\Rightarrow depends (seriously) on δ ?
- YES, but we ignore

For now just open properties ($\delta_i G_{u,i,p}$)

$$\underline{E_{\text{avg}}} \quad \mathcal{F} = \mathcal{F} \Delta \quad ; \quad p_e, \quad \mathcal{S} = \Theta(1/n)$$

(coarse)

more precise: fix λ , $\rho = \lambda/n \rightarrow$

$$E_p [\# \text{ dots}] \sim (np)^3/6 = \lambda^3/6$$

$$\xrightarrow{\approx \text{Poisson}} \mu_p(\mathcal{F}) \approx 1 - \exp[-\lambda^3/6]$$

Ex. 9 $\mathcal{F} = \{ \text{no iso's} \}$ (or $\{ \text{conn's} \}$
 $\{ \exists p, m, e \dots \}$)

~~MD~~ $p = \frac{\ln n + c}{n} \longrightarrow$

$M_p(\mathcal{F}) \longrightarrow e^{-e^{-c}}$
 [EX: nat.]

$\longrightarrow \left\{ \begin{array}{l} p_c \sim \frac{\ln n}{n} \\ \mathcal{F} = O(1/n) \end{array} \right.$
 (Smart p)

Easy: If g is a fixed set of qphs,
then $\exists g := \{ \text{contain some } H \in g \}$
has a coarse threshold.

— is this all?

— wrong question; e.g.:

$\exists = \{ \text{contain } \Delta, \underline{\text{or}} \text{ size} > n \log n \}$
(A) (B)

→ coarse th. (B irrelevant), but ...

usual rough statement of Friedgnt:

if \mathcal{F} has a coarse th. then

~~...~~ \mathcal{F} ~~...~~ \otimes

what's p ? not nec. true at $p_c \rightarrow$

obs: coarse th $\Rightarrow \exists p^* \in \mathcal{G}(\mathcal{F}) \Rightarrow$

$$\frac{d}{dp} \mu_p(\mathcal{F}) \Big|_{p=p^*} < C/p^*$$

~~Thm~~ \otimes holds for $p=p^*$

How it's used:

Thm $\forall \varepsilon, \alpha, C \exists B \exists$

$$p^* \frac{d}{dp} M_p(\mathcal{F}) \Big|_{p^*} < C \frac{1}{B} M_{p^*}(\mathcal{F}) \in (\alpha, 1-\alpha)$$

$\Rightarrow \exists$ (nice) $H \ni$

① $|H| < B$

② $\mu_{p^*}(\exists H) \leq (1/B, 1-1/B)$

③ $B = G_{n, p^*}$, H random copy of H

$$\Rightarrow \Pr(B \cup H \in \mathcal{F}) > 1-\varepsilon$$

$$\left[\dots \exists H \dots \mu_{\mathcal{F}}^*(\mathcal{F}H) \in (1/B, B) \right]$$

Quick Cor If \mathcal{F} has a coarse h_n .

Then $p_{\mathcal{C}}(\mathcal{F}) = \Theta(n^{-\delta})$, some $\delta \in \underline{\mathbb{Q}}$

(e.g.) \rightarrow

$p_{\mathcal{C}}(\mathcal{F}) = \Theta\left(\frac{\log n}{n}\right) \Rightarrow h_n$ is sharp

[monographic:]

original (trickier) appl.:

here usually write $\bar{M} = \#$ of clauses.

(almost*) $\forall k \exists c^* (> 0) \ni \forall \varepsilon > 0$

$$M < (c^* - \varepsilon) n \implies \text{sat. a.s.}$$

$$M > (c^* + \varepsilon) n \implies \underline{\text{unsat. a.s.}}$$

* actually: c^*



OPEN $\lim c_n^*$ exists

Pf of Friedgut (tiny hint):

Fourier transf:

$$f = \mathbb{1}_Z, \quad g = 1 - p \quad (p = p^*) \rightarrow$$

$$\hat{f}(S) = \sum_{A \in \mathcal{F}} \mu_p(A) \left(-\sqrt{\frac{p}{q}} \right)^{|S \cap A|} \sqrt{\frac{q}{p}}^{|S \cap A|}$$

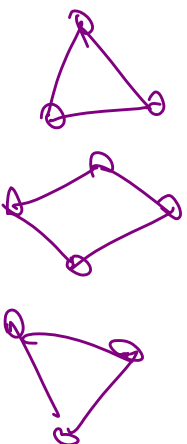
SSX

(just writing,
we don't use)

Very rough = F.T. finds \mathcal{G} :

$\|f\|_2^2$ concentrates on

disjoint U 's of members of \mathcal{G} .



back to derivative

$$\left[p^* \cdot \frac{d}{dp} \mu_p(z) \Big|_{p^*} < C \Rightarrow \dots \right]$$

edge hardy of f :

$$\partial f = \{ (A, A'e) : A \in \mathcal{F}, A'e \notin \mathcal{F} \text{ (} e \in X) \}$$

OBS (Russo, Margulis):

$$p \cdot \frac{d}{dp} \mu_p(z) = p \cdot I_p(z)$$

I_p : "total
influence"

$$:= \sum_{A \in \mathcal{F}} \mu_p(A) \cdot \mathbb{1}_{\{e : A'e \notin \mathcal{F}\}}$$

[meas. of ∂f if meas. of $(A, A'e)$ is $\mu_p(A)$]

edge-isoperimetric inequality:

$$\phi I_p(f) \geq \mu_p(f) \log_p \mu_p(f) \quad (I)$$

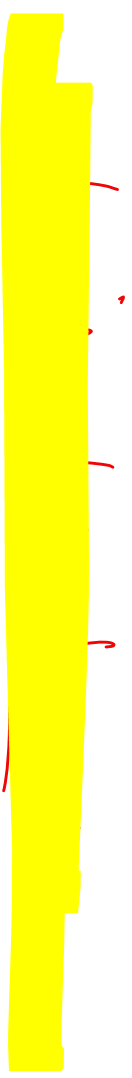
② Sharp: $f = \{ \text{contain } [k] \}$

$(\mu_p(f) = p^k; \text{ each side of } (I) \text{ is } k p^k)$

③ $p \in B(f) \Rightarrow \mu_p(f) = \Omega(1)$

\Rightarrow r.h.s. of (I) = $\Omega(1) \rightarrow$

$\phi = \phi^{k \cdot e}$ $\phi I_p(f) = \phi \cdot \frac{d}{dp} \mu_p(f) < C$



back to KK

$$\left| \Phi_{F_p}(f) \geq \mu_p(f) \log_p \mu_p(f) \right|$$

isoper

Conj 1, $\forall C \exists \varepsilon \exists \delta$

$$\forall f \exists p \in [p_c / \log |X|, p_c] \exists$$

$$\Phi_{F_p}(f) < C \mu_p(f) \log \frac{1}{\mu_p(f)}$$

note missing Φ

Conj 2 $\exists L \exists$

$$\Phi_{F_p}(f) < \mu_p(f) \log \frac{1}{\mu_p(f)} \implies \Phi_E > p/L$$